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$$\log_{10} 2 = 0.30102999566398 \dots$$

x	$\#\{n \leq x : 2^n \text{ begins with } 1\}$
10	3

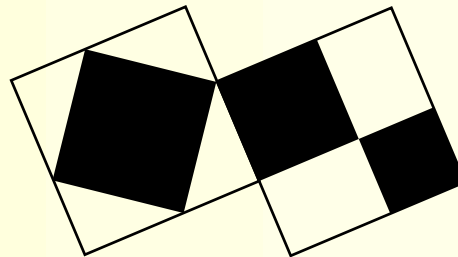
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An Awful Problem about Integers in Base Four
(d'après J H Loxton and A J vdP, *Acta Arith.* 49 (1987), 192–203)



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The matter is troublesome. For instance, given an odd integer k it is not at all obvious how to find a nonzero multiplier m in \mathcal{L} so that also km is in \mathcal{L} . Indeed, the only method we found is not an algorithm at all: it happens always to work, but there's no good a priori reason why it must work.

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$$\begin{array}{r}
 2\bar{1}2\bar{1}1 \quad + \\
 \quad 2\bar{1}2\bar{1}1 \\
 \hline
 11212\bar{1} \quad -
 \end{array}$$

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 110200\bar{1} \\
 \underline{2\bar{1}2\bar{1}1} \quad + \\
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 \underline{2\bar{1}2\bar{1}1} \quad - \\
 110200\bar{1} \quad - \\
 \underline{2\bar{1}2\bar{1}1} \quad + \\
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 \hline
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 \end{array}
 \begin{array}{l}
 + \\
 - \\
 - \\
 + \\
 +
 \end{array}
 \begin{array}{r}
 2\bar{1}11 \\
 \hline
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 \begin{array}{l}
 + \\
 - \\
 - \\
 + \\
 +
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 \begin{array}{r}
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 \underline{2\bar{1}11} \\
 1120\bar{1} \\
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 \underline{2\bar{1}2\bar{1}1} & + & & 0 \\
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 \end{array}$$

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Now denote by \mathcal{S} the set of integers which can be written in base four using just the digits 0 and 1, and for $n = 0, 1, 2, \dots$, denote by \mathcal{S}_n the subset of words in \mathcal{S} of at most n letters.

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$$\mathcal{S}_n + k\mathcal{S}_n = \{s + ks' \mid s, s' \text{ in } \mathcal{S}_n\}$$

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The **so what** of this result is of course that, necessarily, if some element of $\mathcal{S}_n + k\mathcal{S}_n$ has two representatives, say $s_1 + ks'_1 = s_2 + ks'_2$, then

$$k(s'_1 - s'_2) = s_2 - s_1$$

displays a multiplier $s'_1 - s'_2$ in \mathcal{L} yielding $s_2 - s_1$ in \mathcal{L} .

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Given k , say $k \equiv 1 \pmod{4}$, the set $\mathcal{S}_1 + k\mathcal{S}_1$ yields three groups $\{0\}$, $\{1, k\}$, $\{k + 1\}$ consisting of its four elements grouped in congruence classes mod 4.

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Obviously the r_i are bounded in terms of k ; in fact by $(k+1)/3$. Since the r_i must be distinct it follows that **for each k only finitely many different types can occur in the construction.**

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$$\sum_{t \text{ in } \mathcal{S}_n + k\mathcal{S}_n} N_{i-t}^{(n)} = M \quad (0 < i \leq 4^n), \text{ and the given } \sum_{i \pmod{4^n}} N_i^{(n)} = M.$$

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So our attention should turn to the $4^n \times 4^n$ matrix $C = (c_{j-i}^{(n)} - 1)$. It is a **circulant** and those who know such things well well know that it is diagonalisable and that its eigenvalues are given by the 4^n resolvent sums

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The general solution for $N_i^{(n)}$ is given by $4^{-n}M$ from $\theta = 1$ plus some linear combination of solutions coming from the other θ for which $\varphi^{(n)}(\theta)$ vanishes.

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Indeed, it is easy to see by induction that at least 2^n of the $N_i^{(n)}$ are non-zero: for each $i \bmod 4^n$ for which $N_i^{(n)}$ is non-zero, at least two of

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$$M = \sum_{i \bmod 4^n} N_i^{(n)} \geq 2^n.$$

Painting the Lilly

‘To gild refined gold, to paint the lilly, . . . is’, as Salisbury warns King John, ‘wasteful and ridiculous excess’. Nonetheless, we add some remarks on the number of congruence classes of $\mathcal{S} + k\mathcal{S} \bmod 4^n$, and therefore an alternate proof, primarily, I guess, because that was our original line of argument.

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Thus if we follow a type of say h elements to level n we either find fewer than $4^n h$ elements or we find a congruence class mod 4^n with more than h elements. In either case at least one of the 4^n congruence classes must contain fewer than h elements. **So each type leads eventually to the singleton type.**

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Incidentally, the argument fails if the last nonzero digit of k is a 2 , because T then has an irreducible component in which all row sums are 4 .

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It follows that if almost all the $r_n(k)$ are zero then some $r_n(k)$ must exceed 1, again solving our problem. Our arguments in fact show, if k is odd, that there are $r_n(k)$ that are arbitrarily large.

Notes and References

Gavin Brown, William Moran, and Robert Tijdeman, 'Riesz products are basic measures', *J. London Math. Soc.* **30** (1984), 105–109.

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D. H. Lehmer, K. Mahler and A. J. vdP, 'Integers with digits 0 or 1', *Math. Comp.* **46** (1986), 683–689.

We knew that $\mathcal{S} - \mathcal{S} = \mathcal{L}$ because of this work.

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This talk, though without my spoken commentary, can be found at <http://www.maths.mq.edu.au/~alf/AwfulTalk.pdf>.

Gavin

$$V \times XIII = XIII \times V$$

Gavin

Happy

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Gavin

Many Happy

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Many Happy Returns

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